



Security Holdings

BLC: Blocklist Co-occurrence Analysis for Large-scale IP Network Traffic Flows

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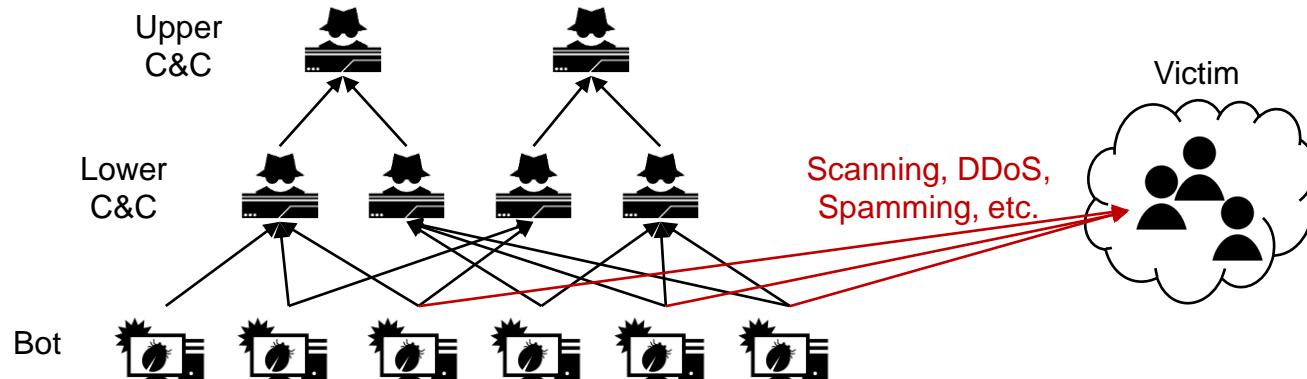
Background



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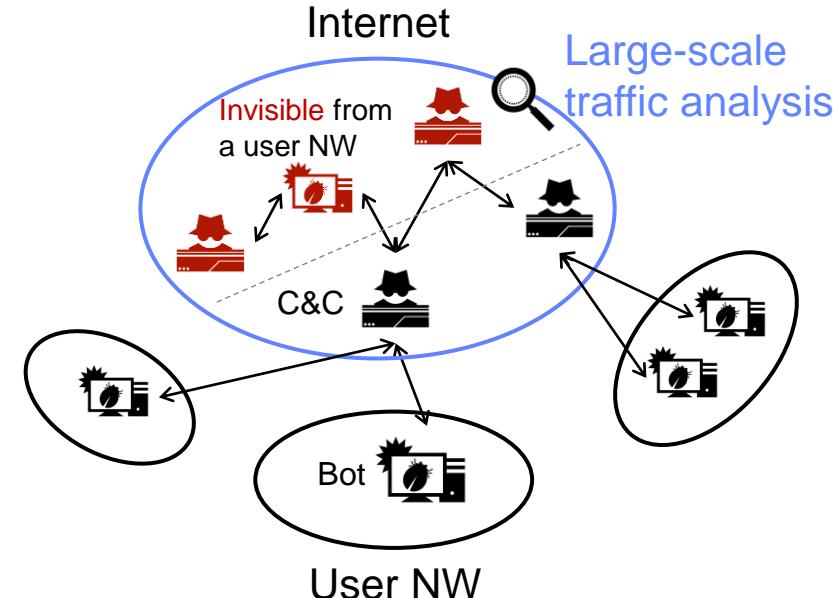
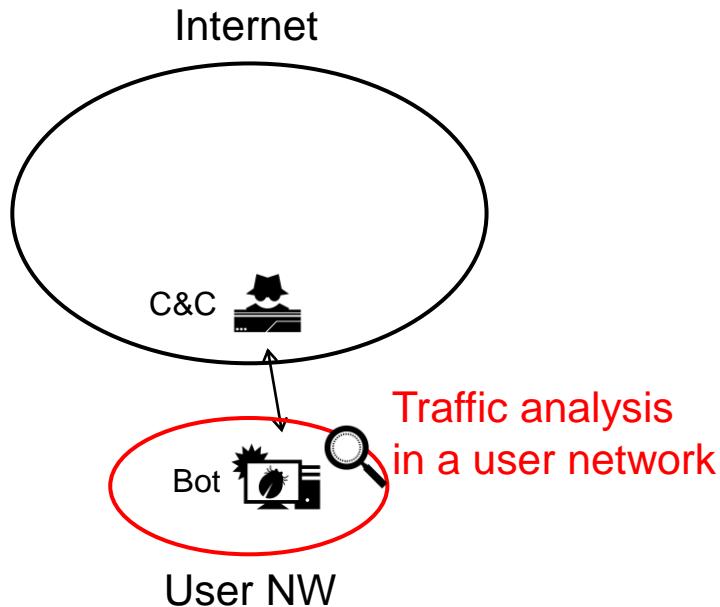
- A botnet is a group of malware-infected hosts that launch various cyberattacks.
- To eliminate botnet threats, as the first step we need to understand the entire picture of botnet infrastructures by
 - detecting components in a botnet, and
 - identifying the relationships between them

Example of hierarchical botnet structure



Problem Statement 1/3

- Many existing methods focus on traffic analysis in a user network.
- However, they lack the visibility of layered and distributed botnet infrastructure.
- Large-scale traffic analysis at the Internet backbone is necessary.



Problem Statement 2/3

C&C detection by ML

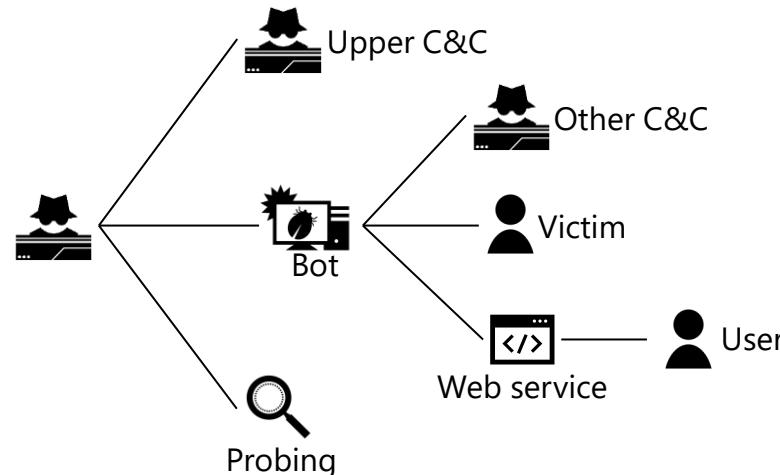
Detect C&C by analyzing their communication behaviors individually but not collectively



- Quantity
- Frequency
- Number of contacts
- Time of day

Graph approach

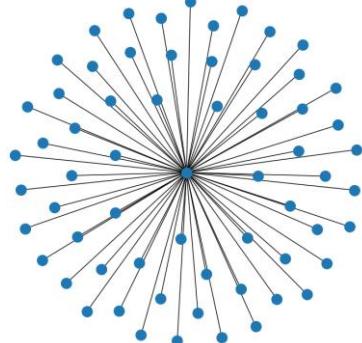
Grasp a group (or botnet) of collaborating hosts by analyzing communication between them



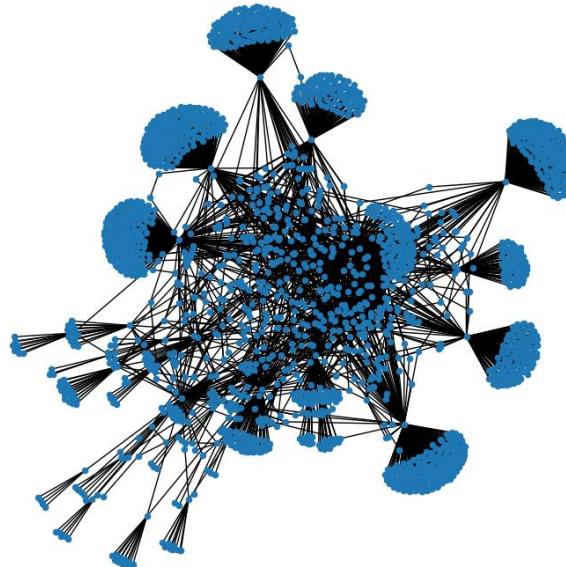
Problem Statement 3/3

- In an Internet-scale network, the number of neighbors of a given host can be very large.
- Therefore, it is essential to have an algorithm that can **efficiently** detect malicious hosts from the **huge graph**.

1-hop neighborhood

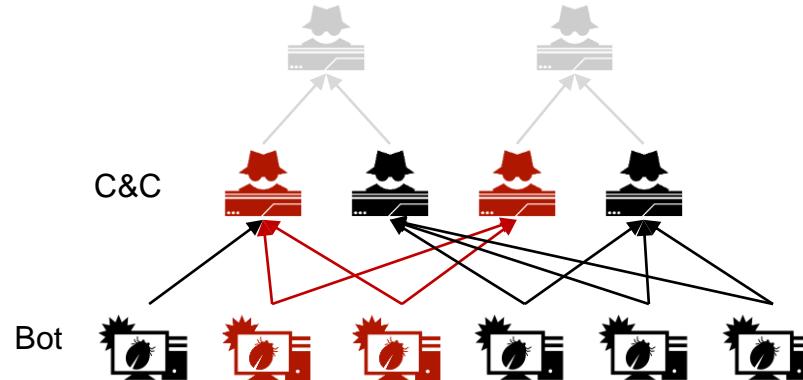


2-hop neighborhood



Basic Idea for Detection

- We propose a novel detection method, **BLC** (**B**lock**L**ist **C**o-occurrence analysis).
- To detect malicious servers on the IP graph, we use **detection by co-occurrence** [1].
 - Malicious actors might prepare not only 1 malicious server but several servers.
 - Infected hosts might connect to several malicious servers.
- For efficient detection on huge graph, we additionally use **pruning technique**.



[1] K. Sato, K. Ishibashi, T. Toyono, H. Hasegawa, and H. Yoshino, “Extending black domain name list by using co-occurrence relation between DNS queries,” IEICE Trans. Commun., vol. E95.B, no. 3, pp. 794–802, Mar. 2012.

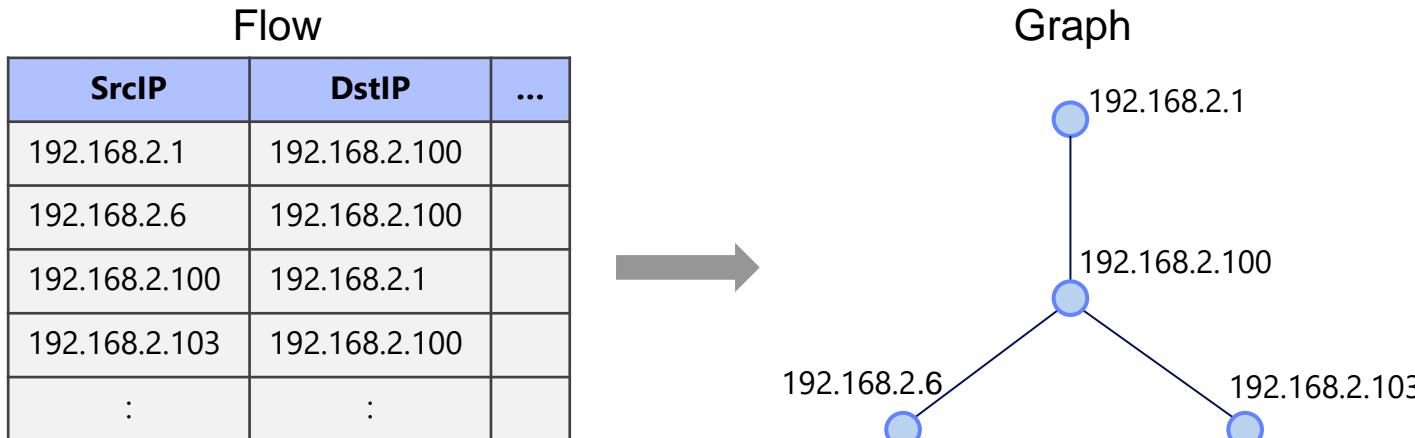
Proposal: Graph Construction



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- We generate a graph of communication relationships between IP addresses from flow data.
- The graph is undirected, because the sampled flow data do not always tell us which host initiated the communication.

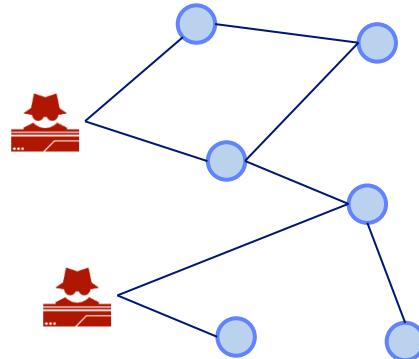


Proposal: Listing Up Candidates

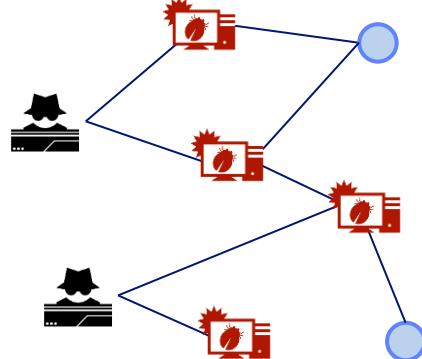


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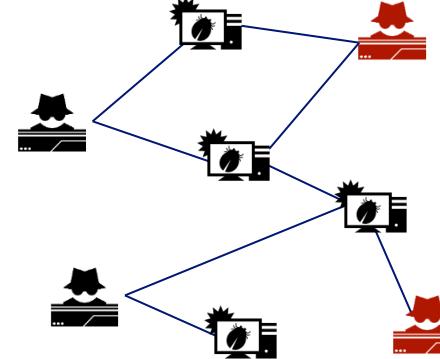
- We assume all IPs communicating with blocklist IPs are bots.
- IPs communicating with bots that are neither bots nor blocklist IPs are considered malicious IP candidates.



Seeds: Blocklist



1-hop: Bot Candidates



2-hop: C&C Candidates

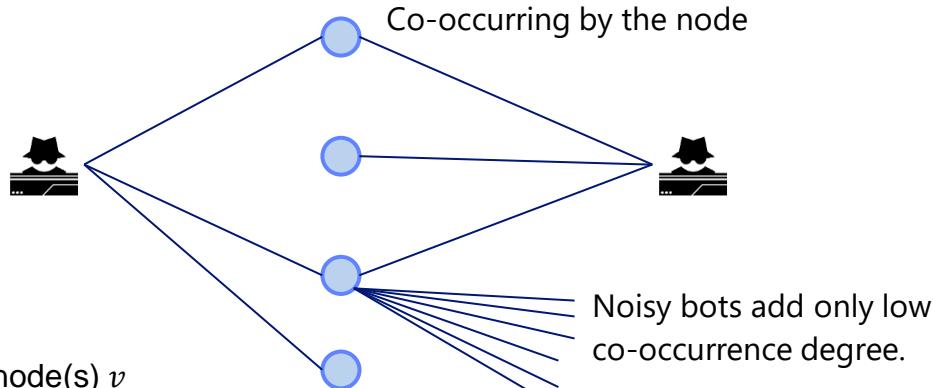
Proposal: Score Calculation 1/2



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- Co-occurrence between hosts is calculated by the similarity of the communication destination set.
- To reduce the effect of noisy nodes such as scanners, it is calculated as a **weighted Jaccard Index** of neighborhood as follows:

$$C(h_i, h_j) = \frac{\sum_{v \in N(h_i) \cap N(h_j)} 1/|N(v)|}{|N(h_i) \cup N(h_j)|}$$



* $N(v)$: neighborhood of node(s) v
 h_i : i -th host

Proposal: Score Calculation 2/2

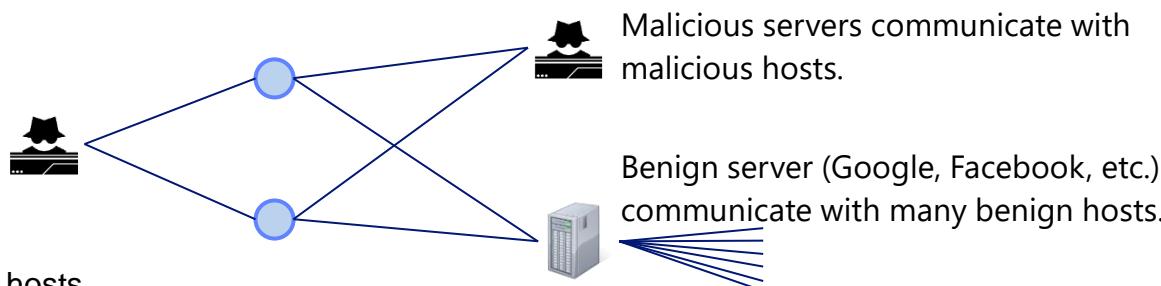


- Since bots often communicate with benign services to check for connection, we consider malicious weight by **rate of bots in neighborhood**:

$$W(h_i) = \frac{|N(H_{mal}) \cap N(h_i)|}{|N(h_i)|}$$

- The final malicious score is **product of co-occurrence and malicious weight**:

$$\text{Mal}(h_i) = W(h_i) \times \sum_{h \in H_{mal}} C(h, h_i)$$



* H_{mal} : Set of blocklist hosts

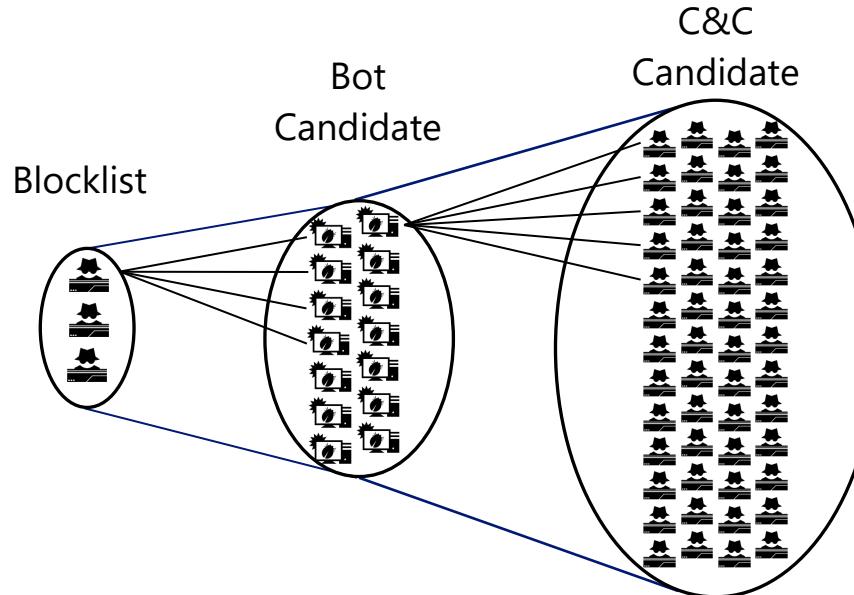
Proposal: Pruning Technique 1/2



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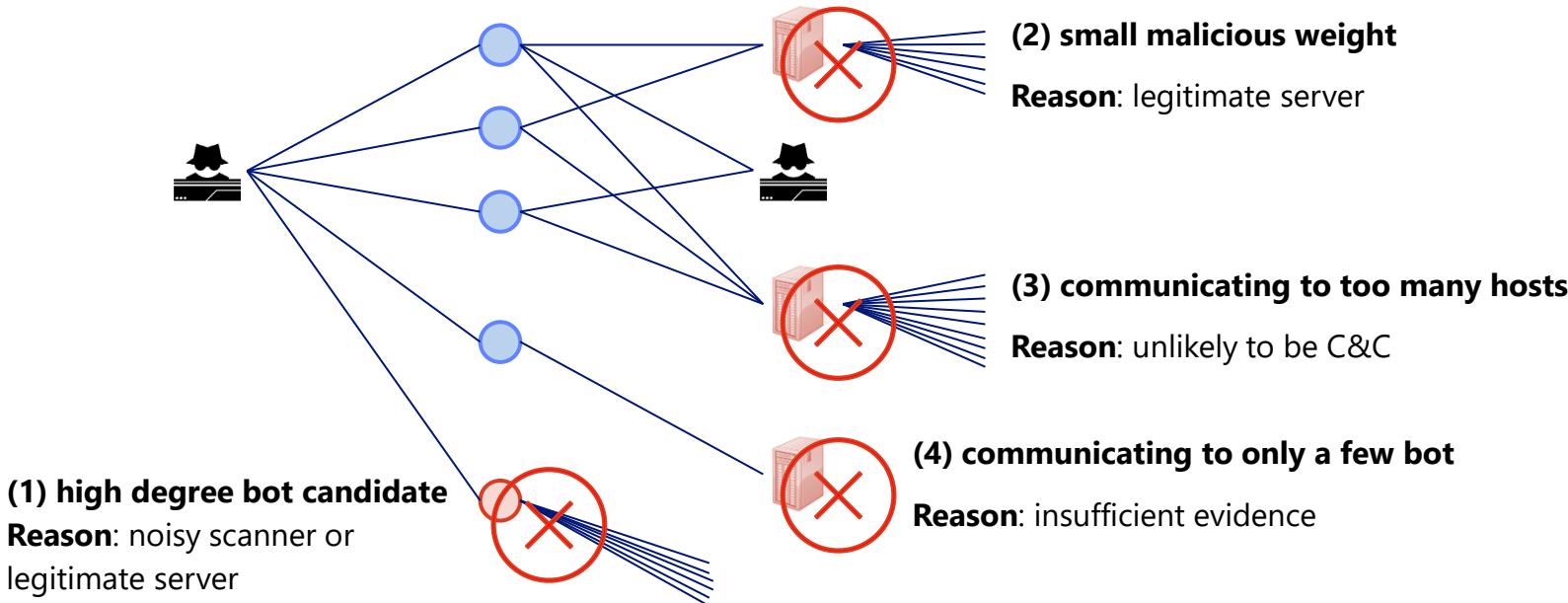
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- To calculate the scores for all C&C candidates, we need the weighted Jaccard index calculation of (# block list IPs) x (#C&C candidate IPs) times.
- This is a very time-consuming operation because the size of a 2-hop neighborhood can be very large in an Internet-scale graph.



Proposal: Pruning Technique 2/2

- The following pruning heuristic is applied.



Evaluation: Dataset and Parameters

- Flow data: Real flow data of a large network
- Blocklist: Seed blocklist (general C&C, IoT C&C)

Statistics of data:

Item	Size
Flow records (per day)	2.2×10^9
Graph nodes (per day)	1.8×10^8
Graph edges (per day)	6.3×10^8
General C&C blocklist	3932
IoT C&C blocklist	483

proposed pruning parameter:

Parameter	Meaning
bot_deg = 100	degree of bot candidate > bot_deg
weight = 0.1	malicious weight > weight
c2_deg = 1000	degree of C&C candidate < c2_deg
comm_bot = 4	# of communicating bots > comm_bot

Evaluation: Validation with Analyst 1/2

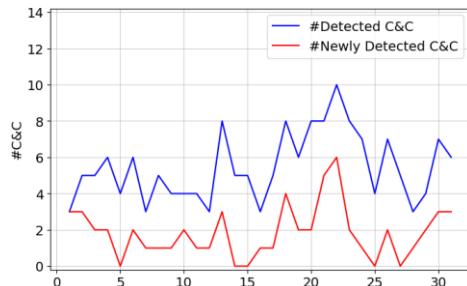
The result of the **proposed method** (BLC)

- We extract 100 most suspicious host per day for 1 month.
- The result is validated by a security analyst using OSINT, which is independent of the seed block list.

Used blocklist

General C&C
blocklist

#detected C&C server



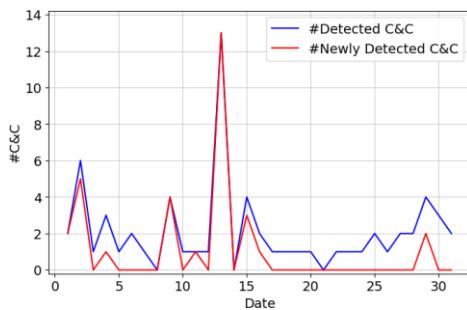
#C&C / #Unique IPs

57 / 909

**Percentage expansion
of block list**

57 / 3932
= 1.4%

IoT C&C
blocklist



32 / 862

32 / 483
= 6.6%

Evaluation: Validation with Analyst 2/2

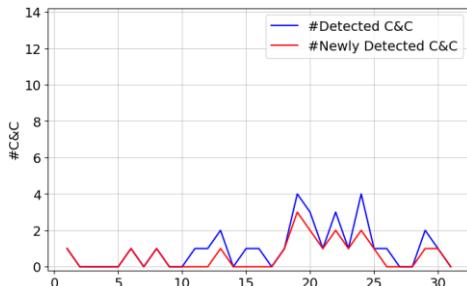
The result of the **conventional method** (BLC without pruning)

- It **detects fewer C&C servers** than the proposed method.

Used blocklist

General C&C
blocklist

#detected C&C server



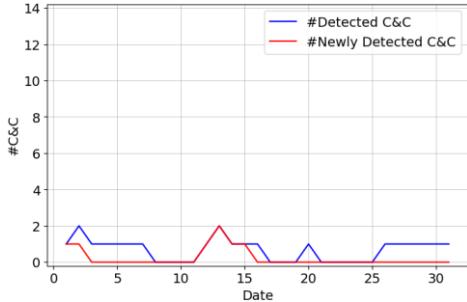
#C&C / #Unique IPs

19 / 1382

Percentage expansion of block list

19 / 3932
= 0.5%

IoT C&C
blocklist



7 / 859

7 / 483
= 1.5%

Evaluation: Effect of Pruning

Change in computation time with pruning

- Proposed pruning significantly reduces computation time.
- The score calculation is parallelized on 32 cores.

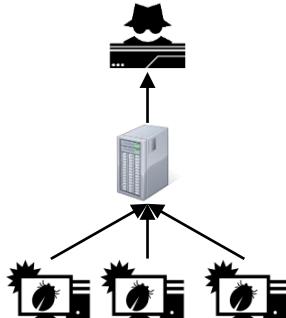
Method	#2-hop nodes	Process time (minutes)	Similarity* of top 100 IPs with conventional method
Conventional method [1] (BLC without pruning)	5.4×10^7	161.4	100%
BLC (bot_deg=3000, weight=0.1)	8.9×10^6	26.6	98%
BLC (proposed pruning)	1.4×10^5	0.7	28%

*by Jaccard index

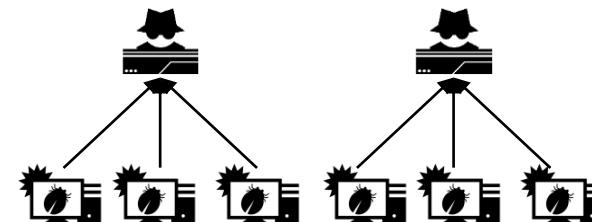
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- The proposed method assumes that bots communicate directly with multiple C&C servers. Therefore, It is not effective in the following cases:

C&C and bots communicate through proxy servers.



Each bot communicates with only one C&C.
(C&C servers are not connected.)



Conclusion

- We propose a method (BLC) to detect malicious hosts related to the given blocklist hosts from flow data.
- It works efficiently even for Internet-scale IP network traffic flows and is more than 100 times faster and higher precision than conventional method.
- Evaluation using large real flow data show that BLC find many C&C servers.

Appendix: Malware Types



- BLC detects hosts that are the C&C servers for the following malwares.

from general C&C blocklist

Malname	Type	#IP
win.socks5_systemz	Other	35
elf.mirai	DDoS	6
elf.mozi	DDoS	3
unknown.unknown	Other	3
win.asyncrat	RAT	1
script.coinMiner	Unknown	1
elf.bashlite	DDoS	1
win.rhadamanthys	CredentialStealer	1
win.teamspy	RAT	1
8220-Gang	Unknown	1
win.cobalt_strike	PentestFramework	1
elf.gafgyt	DDoS	1
dll	Unknown	1
elf.unknown	Other	1

from IoT C&C blocklist

Malname	Type	#IP
elf.mirai	DDoS	8
elf.bashlite	DDoS	4
win.cobalt_strike	PentestFramework	3
script.unknown	Other	2
win.dcrat	RAT	2
win.redline_stealer	CredentialStealer	1
win.nanocore	RAT	1
win.quasar_rat	RAT	1
elf.unknown	Other	1
win.shadowpad	Backdoor	1
win.icedid_downloader	CredentialStealer	1
win.sliver	PentestFramework	1
ascii.unknown	Other	1
elf.opendir	Other	1
win.bazarbackdoor	Backdoor	1
script.coinMiner	Unknown	1
win.icedid	CredentialStealer	1